



Protecting Privacy Against Structural De-anonymization Attacks in Social Networks

Public defense of

Gábor György Gulyás

Supervisor: *Sándor Imre, DSc*

Reviewers:

Gergely Biczók, PhD (BME-TMIT)

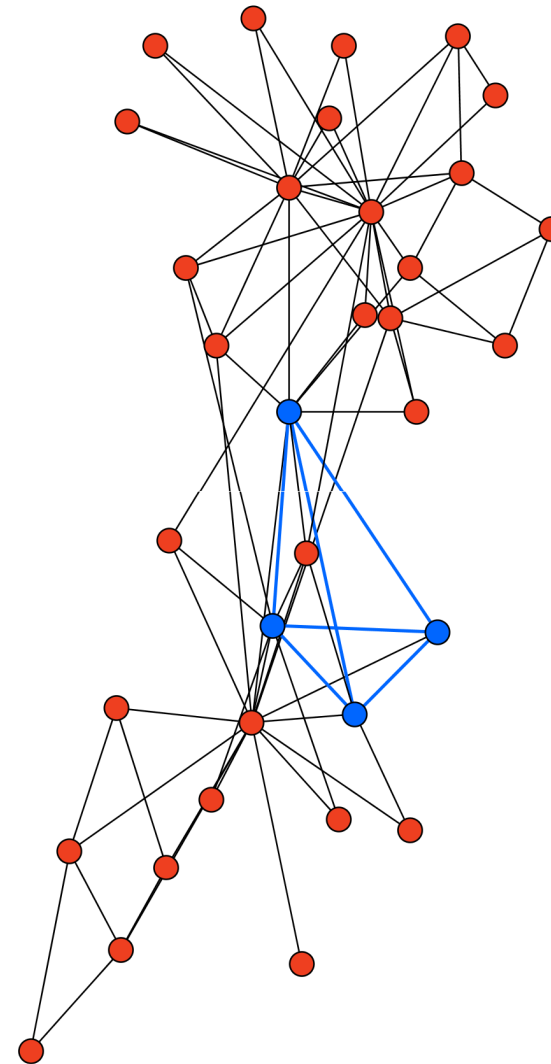
Julien Freudiger, PhD (PARC, USA)



MOTIVATION

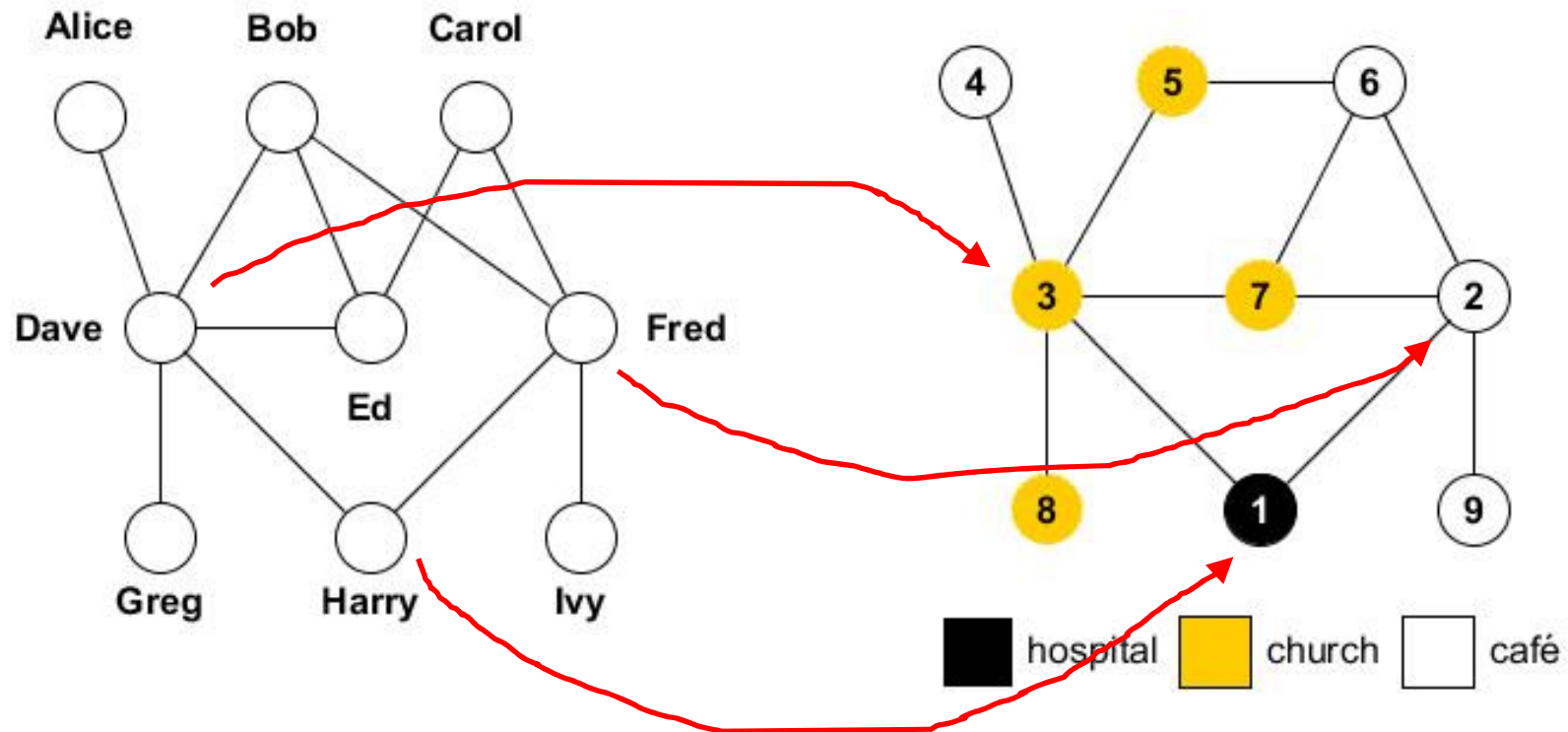
Structural privacy in social networks

- One of most challenging privacy issues: how structural information can be hidden
- Can be leaked [anonymously] easily
 - TP businesses
 - TP research partners
 - Implicit social networks
 - E-mail networks
 - Call information
 - ...



Structural privacy in social networks (2)

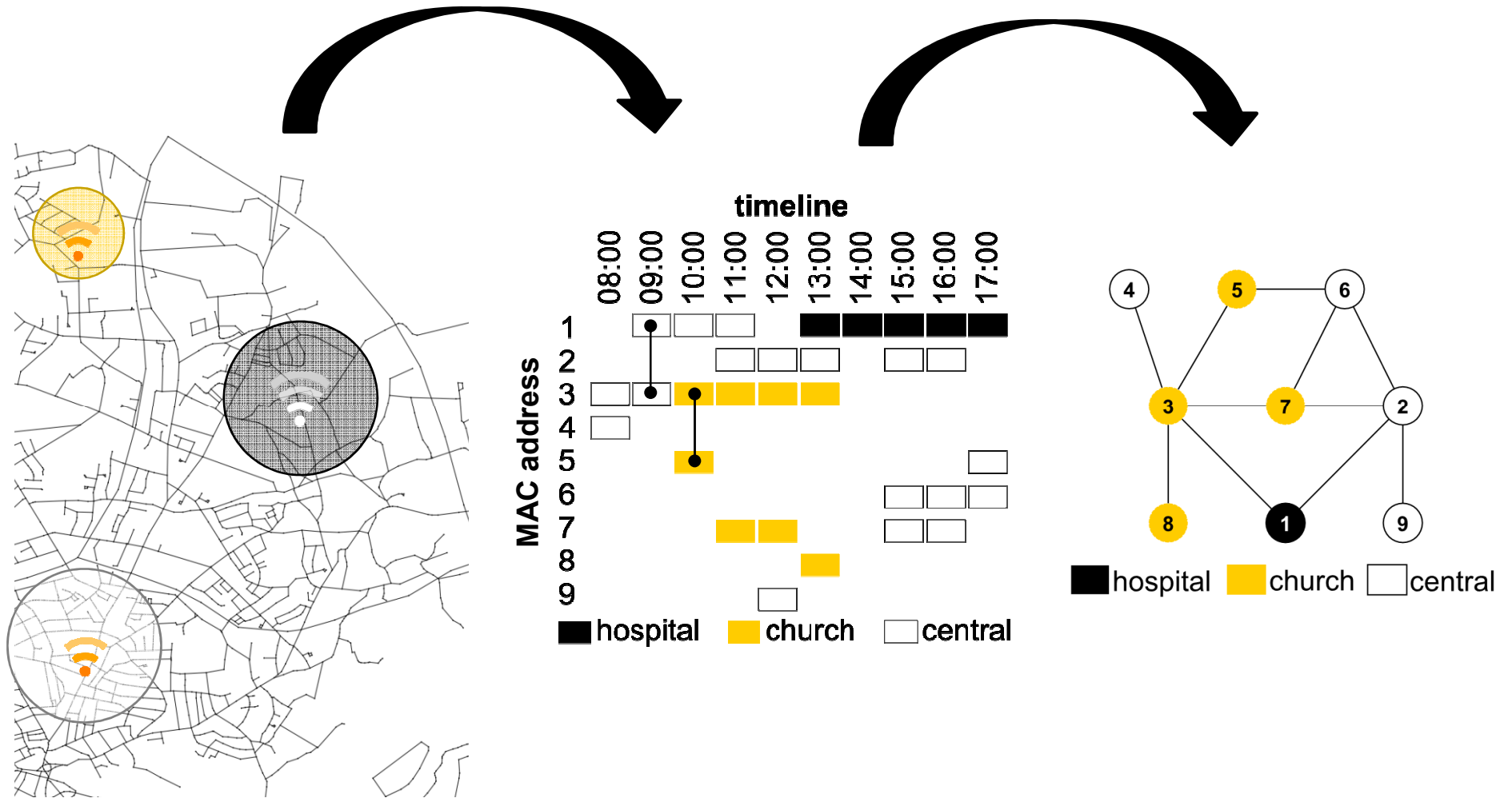
Phases:
1. Initialization (seeding)
2. Propagation



Goals:

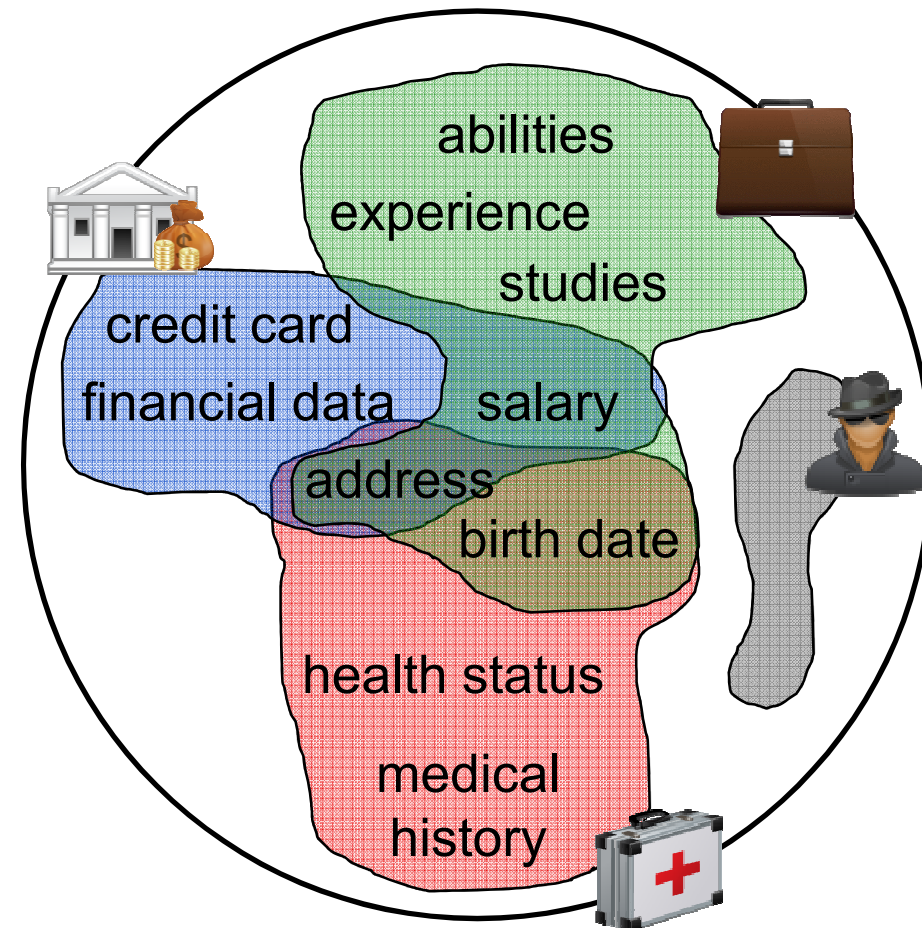
- Attacker: achieve high recall
- Users: minimize network or individual information disclosure

Structural privacy in social networks (3)



How to protect user privacy?

- Proposed solution:
 - Identity separation
- Desired properties?
 - Work with existing services, gradually adoptable
 - User control: client side solutions
 - Without consent of any service providers
 - All actions on client side



**Global and partial identities
of John Doe**



NEW RESULTS

Problem sets of the dissertation

Analysis of re-identification attacks

- Measuring anonymity
- Analysis of the initialization of attacks

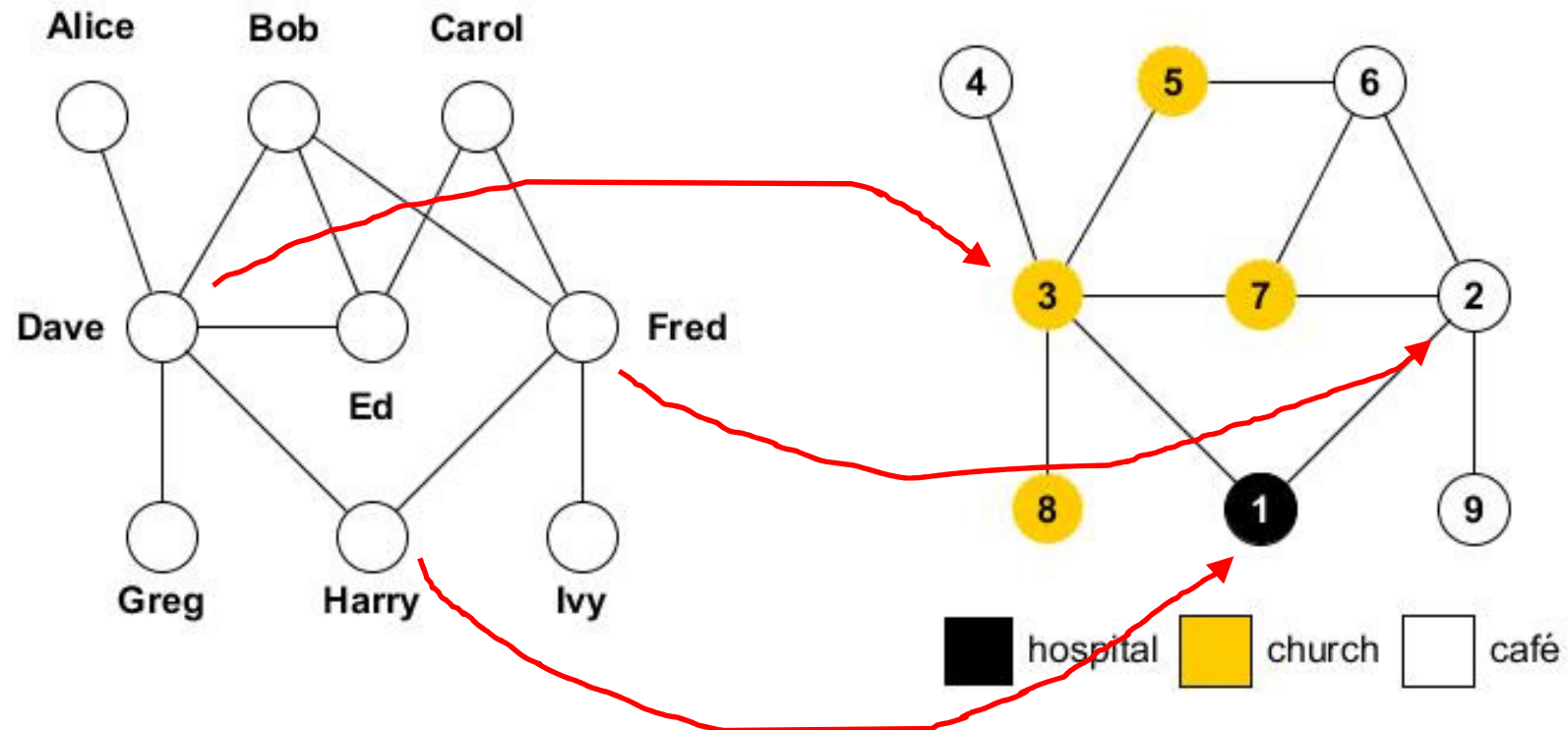
Evaluation of identity separation

- Against attack initialization
- Evaluation of strategies to stop re-identification
 - Individual adoption
 - Cooperation

Individual strategies with identity separation

- Small adoption rates and use of decoys
- Reverseability
- K-anonymity
- Game-theoretic approach

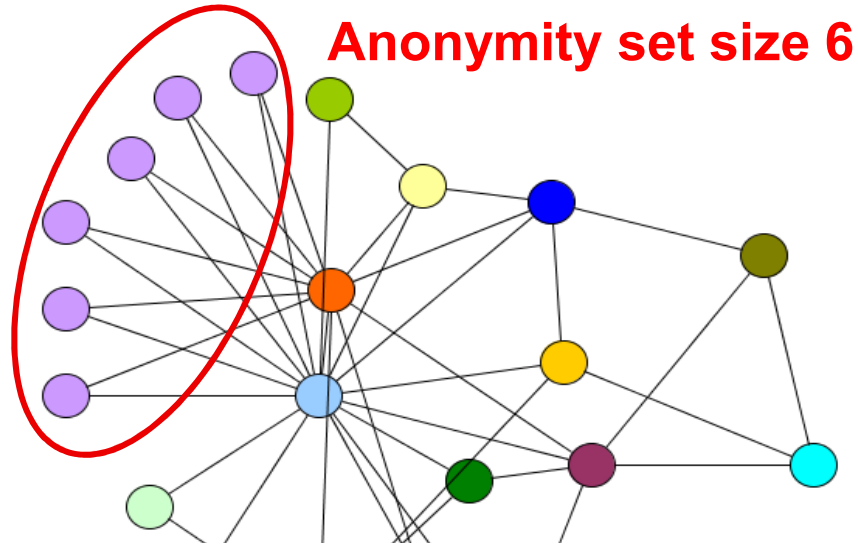
Problem Set 1: *analysis of re-identification attacks*



a. Measuring anonymity?
– No global anonymity sets

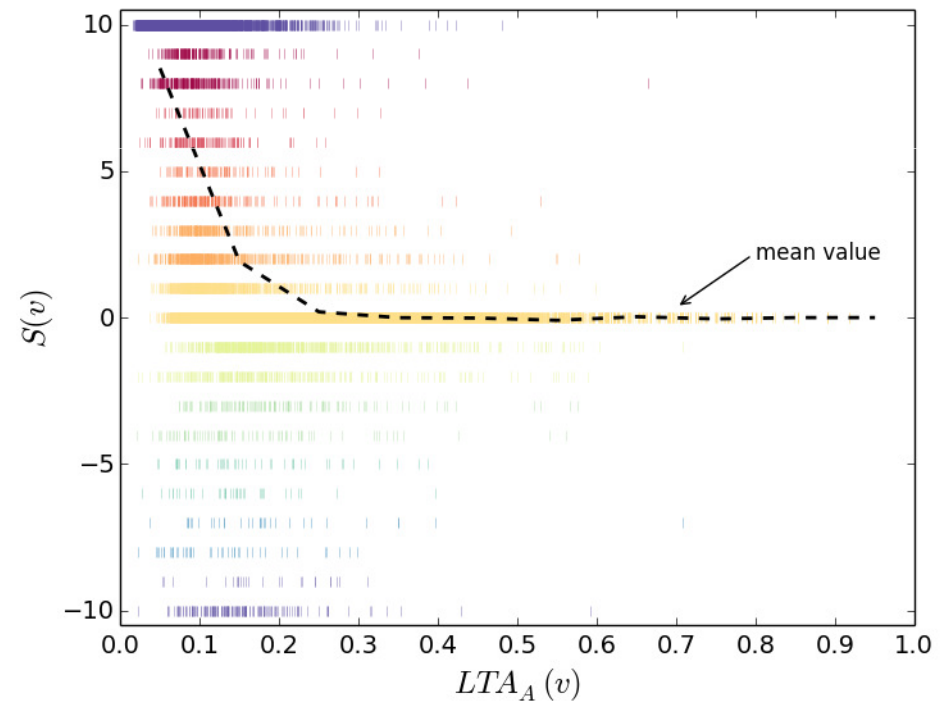
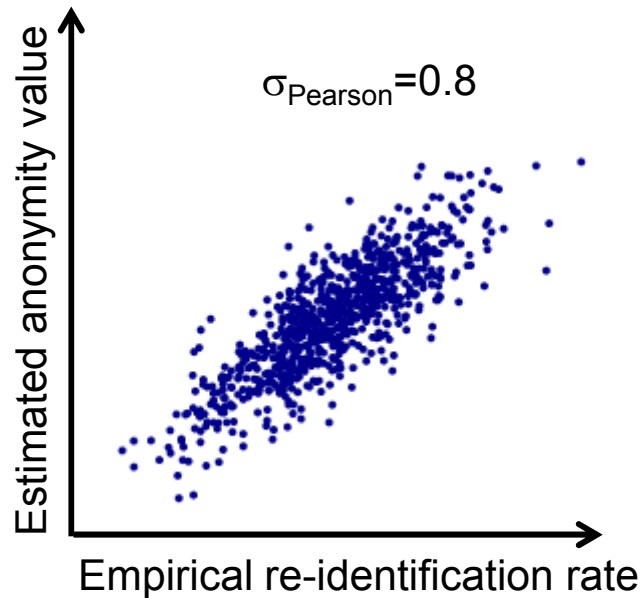
b. 2-phase attacks:
– seeding (initialization)
– propagation

[T1.1-2] Can we measure anonymity?

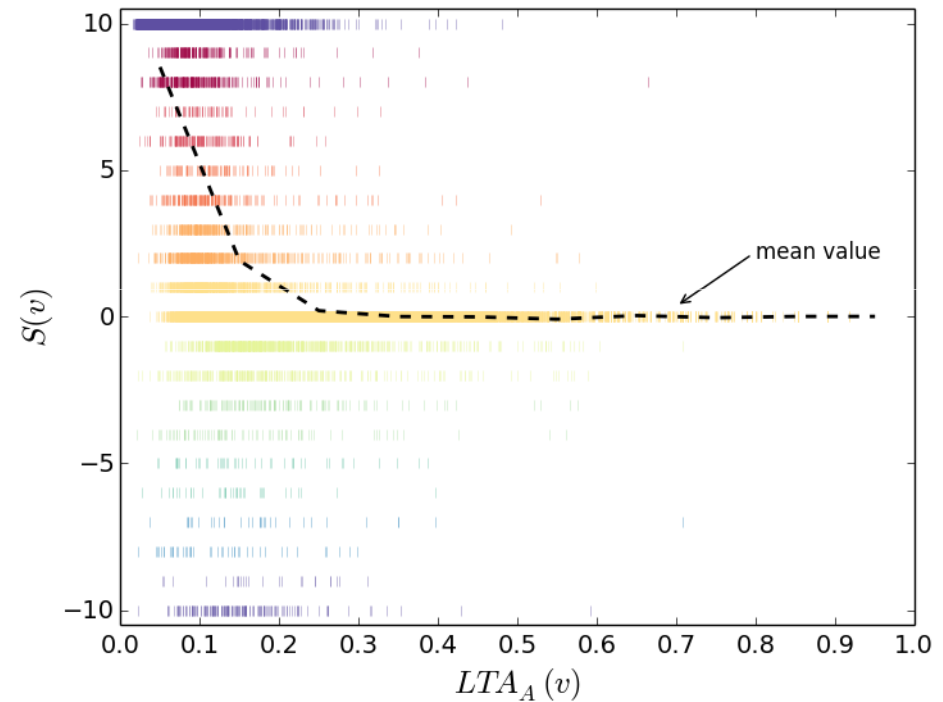
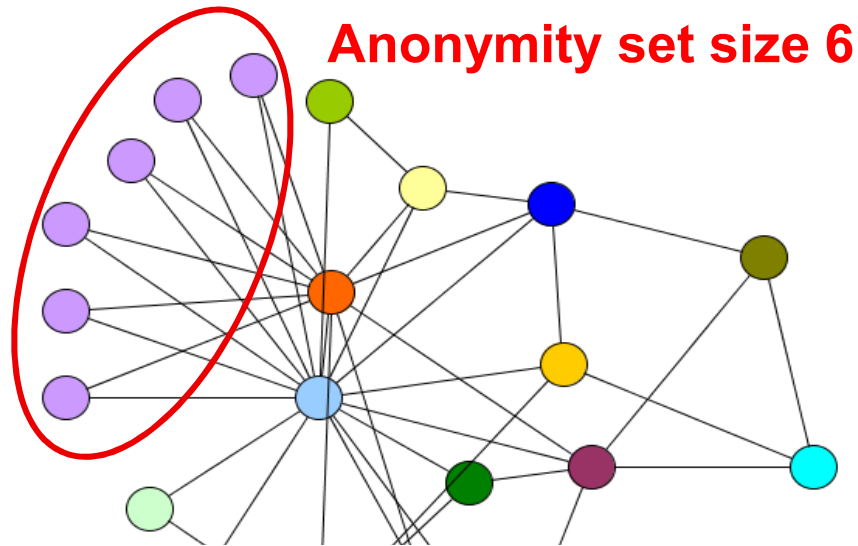


$$LTA_{deg}(v_i) = \text{deg}(v_i).$$

$$LTA_A(v_i) = \sum_{\forall v_k \in V_i^2} \frac{\text{CosSim}(v_i, v_k)}{|V_i^2|},$$



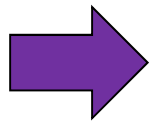
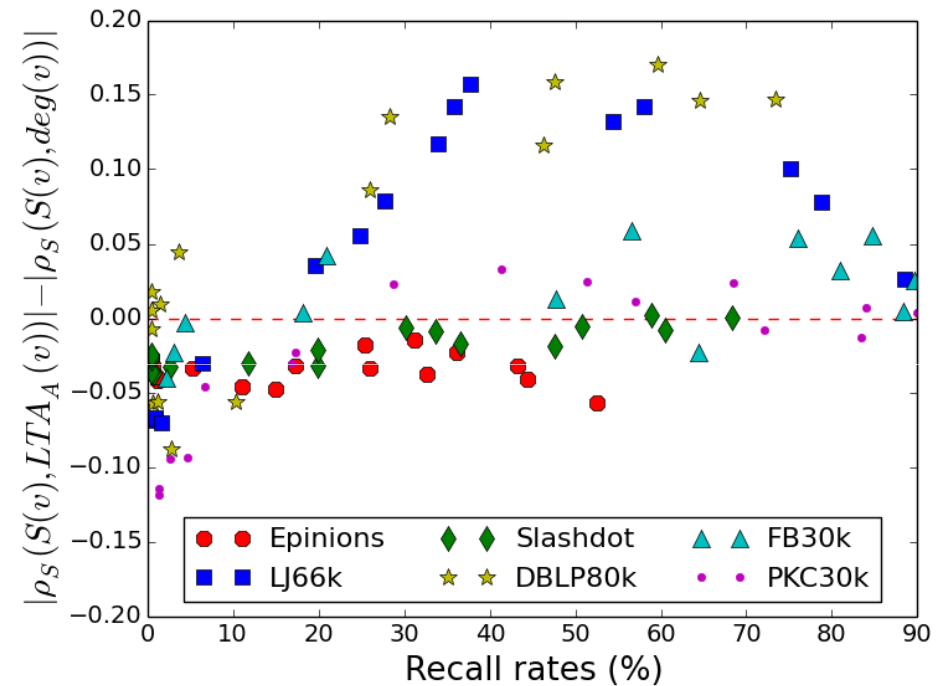
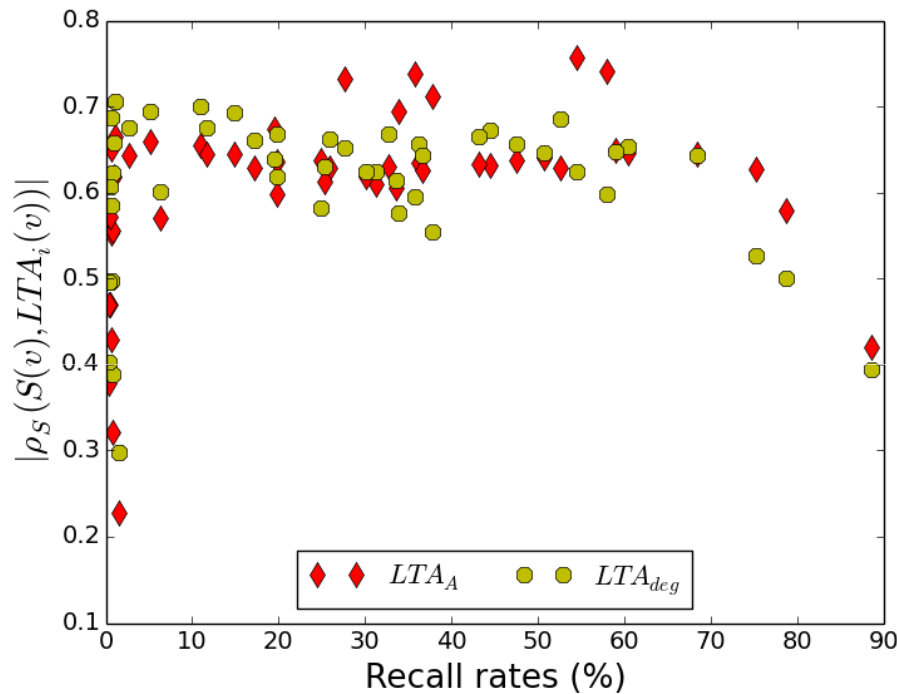
[T1.1-2] Can we measure anonymity?



$$LTA_{deg}(v_i) = \deg(v_i).$$

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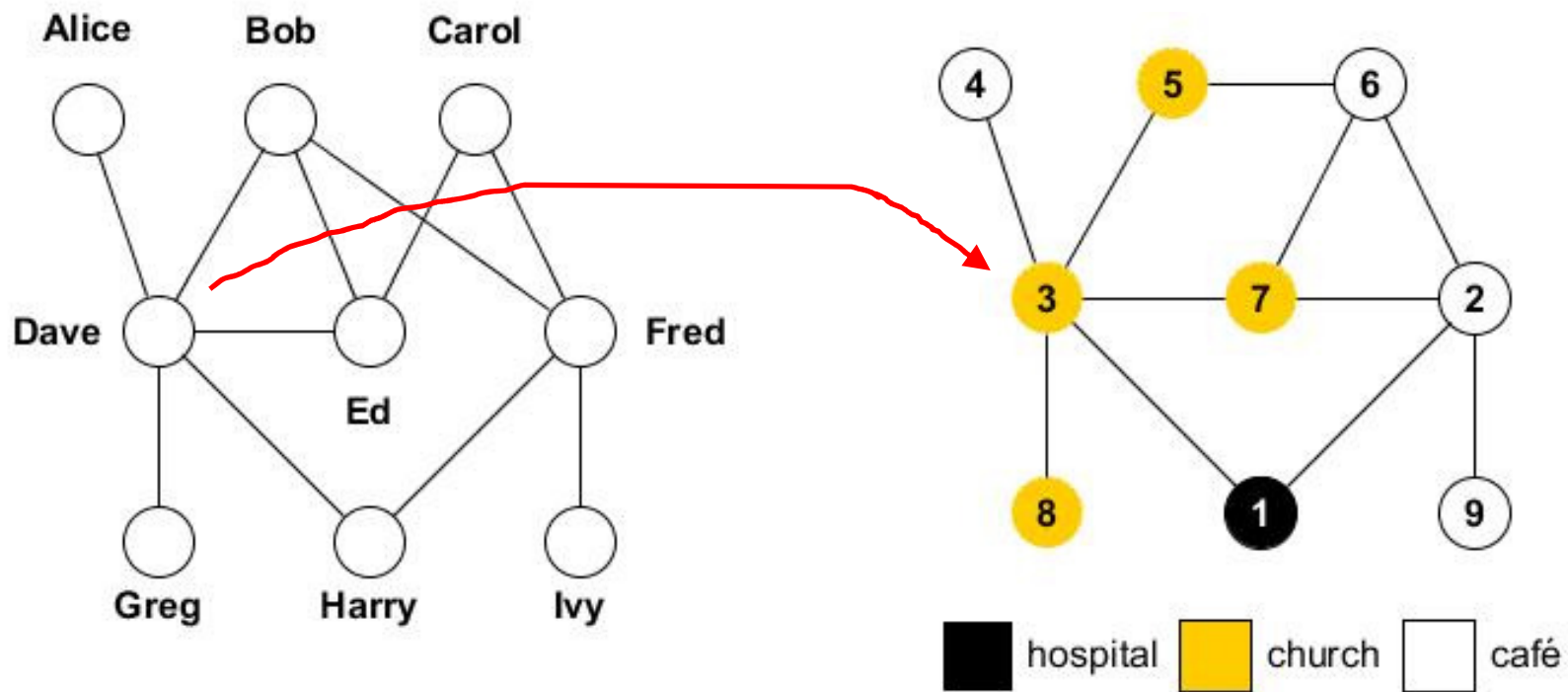
[T1.1-2] Can we measure anonymity? (2)



Applications, e.g.:

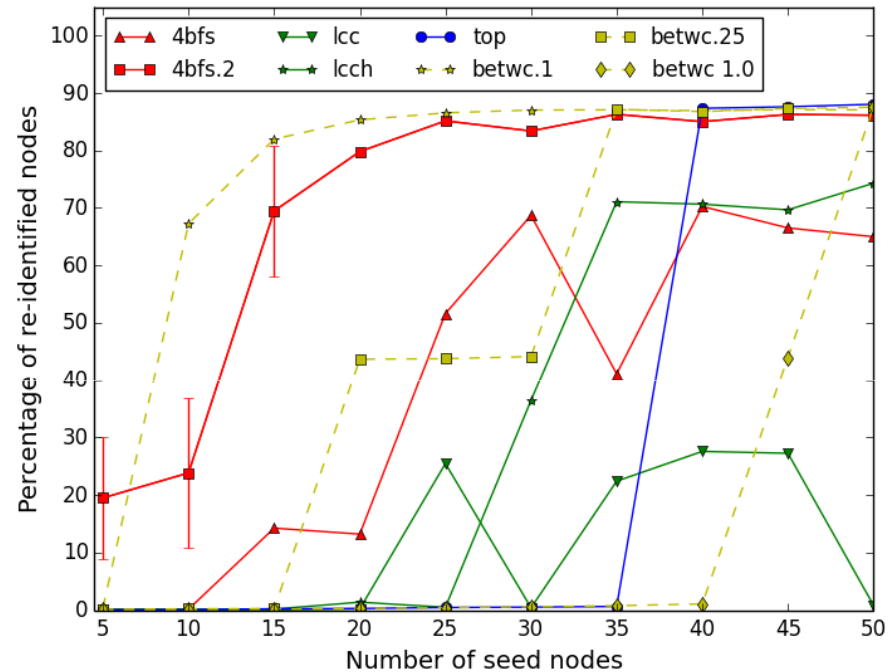
- individual / provider risk assessment
- as measures of importance

[T1.3] How initialization effects overall results?



[T1.3] How initialization effects overall results? (2)

- For the state-of-the-art algorithm, I characterized the importance of initialization.
- In particular:
 - I showed how the maximum number of re-identified nodes can depend on the seeding method and its parameters
 - I showed how the minimum number of seed nodes depends on network properties and the seeding method
 - I characterized seed stability and showed that even an extremely low number of seed nodes can also lead to large-scale propagation



Seeding method

- should be carefully chosen
- part of the attacker model

Related publications to Problem Set 1

[J2] B. Simon, G. G. Gulyás, and S. Imre, “Analysis of grasshopper, a novel social network de-anonymization algorithm,” *Periodica Polytechnica Electrical Engineering and Computer Science*, January 2015. (accepted for publication).

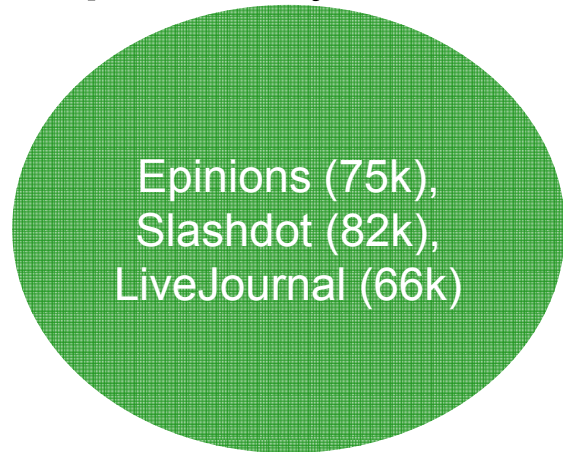
[J3] G. G. Gulyás and S. Imre, “Using identity separation against de-anonymization of social networks,” *Transactions on Data Privacy*, January 2015. (accepted for publication).

[C1] G. G. Gulyás and S. Imre, “Measuring importance of seeding for structural de-anonymization attacks in social networks,” in *Pervasive Computing and Communications Workshops (PERCOM Workshops), 2014 IEEE International Conference on*, 2014.

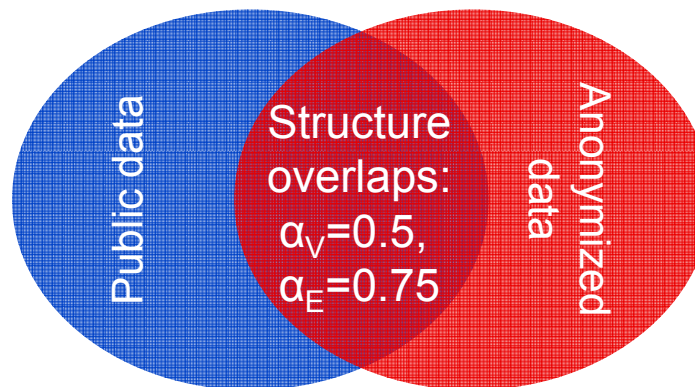
[C3] G. G. Gulyás and S. Imre, “Measuring local topological anonymity in social networks,” in *Data Mining Workshops (ICDMW), 2012 IEEE 12th International Conference on*, pp. 563–570, 2012.

Problem Set 2: *evaluation of identity separation*

Step 1: anonymized network



Step 2: perturbation

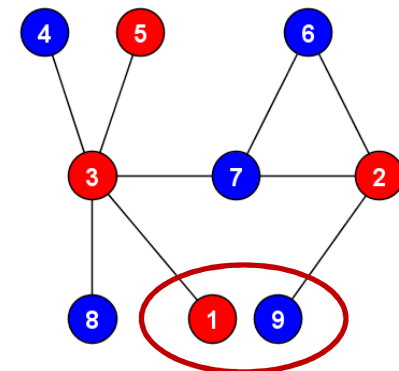


ground truth

Multiple models

- Number of identities (Y)
- Deleting edges, duplicating edges

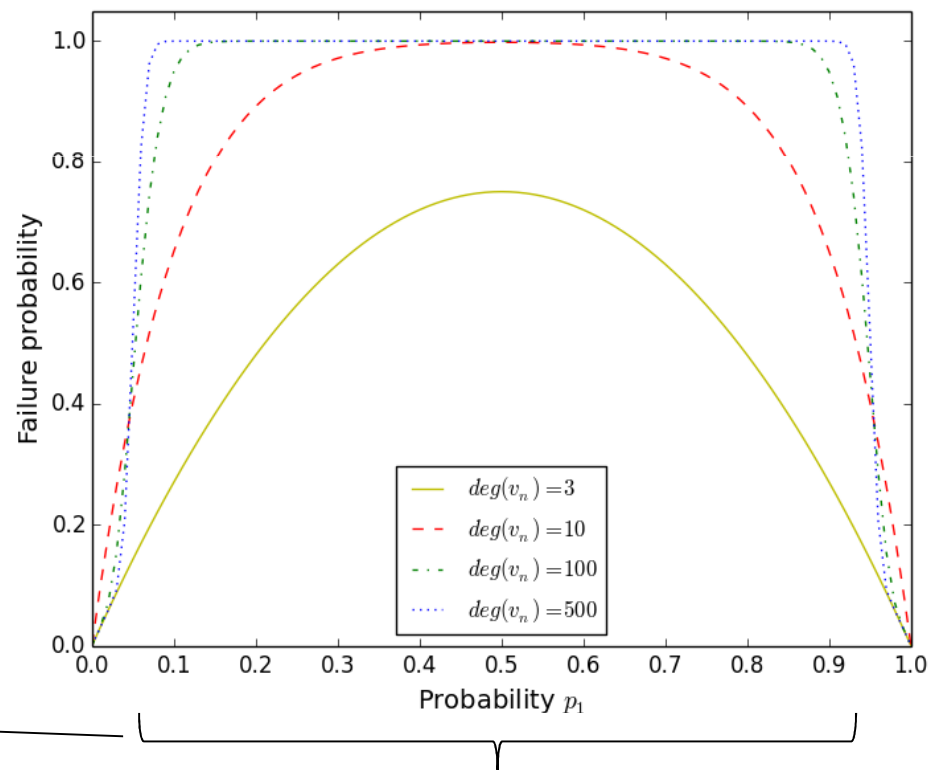
Step 3: simulating identity separation



[T2.1] Can we estimate failure probability of seeding?

By providing the general formula of failure probability, I elaborated the lower estimate of failure probability for clique-based, and top node based seeding:

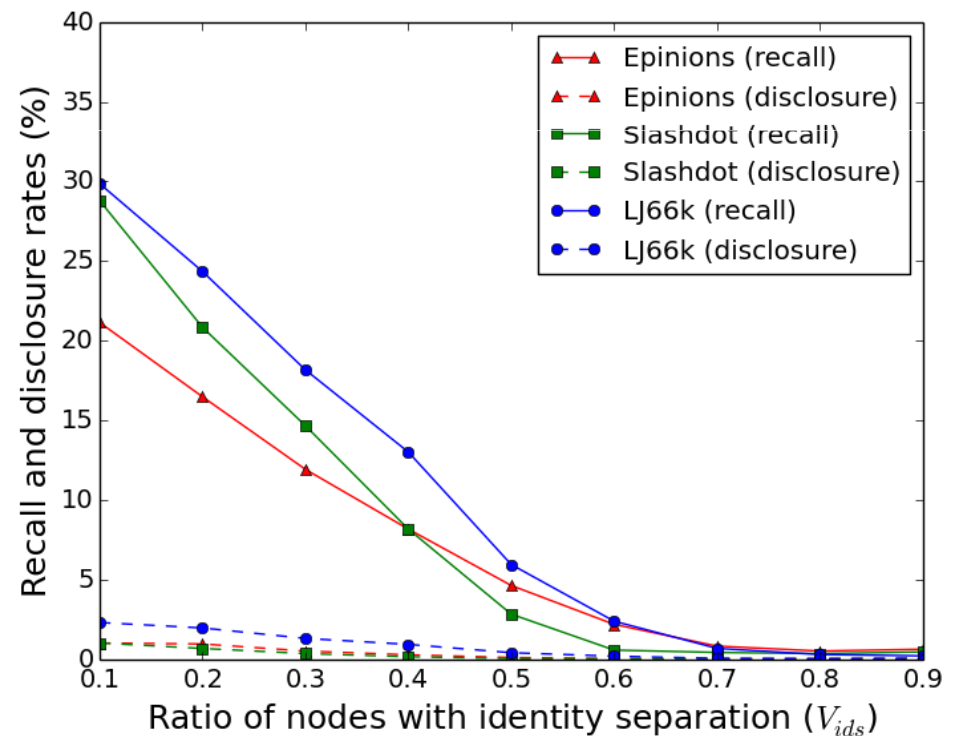
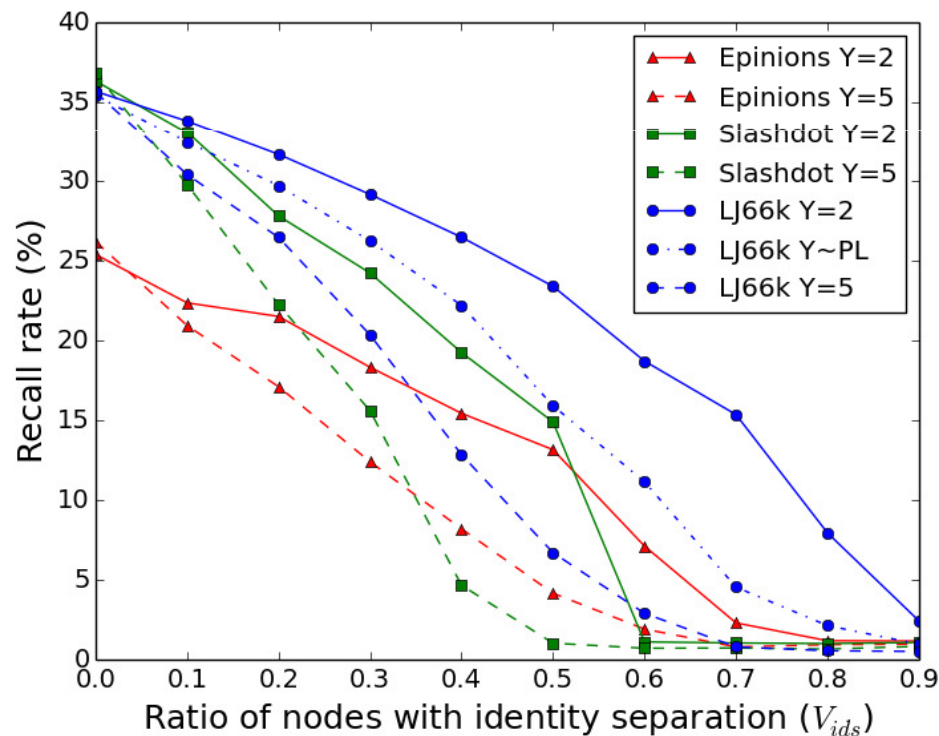
$$P_{clique}^B(\text{"failure"}|Y = y) = 1 + \sum_{\forall i \in [0, \dots, y]} p_i^{k-1} \cdot \left(\sum_{x_1'' + \dots + x_y'' = n-k+1} \left(\frac{(n-k+1)!}{x_1''! \cdot \dots \cdot x_y''!} \cdot p_1^{x_1''} \cdot \dots \cdot p_y^{x_y''} \cdot e(k-1, x_i'') \right) - 1 \right) \quad (3)$$



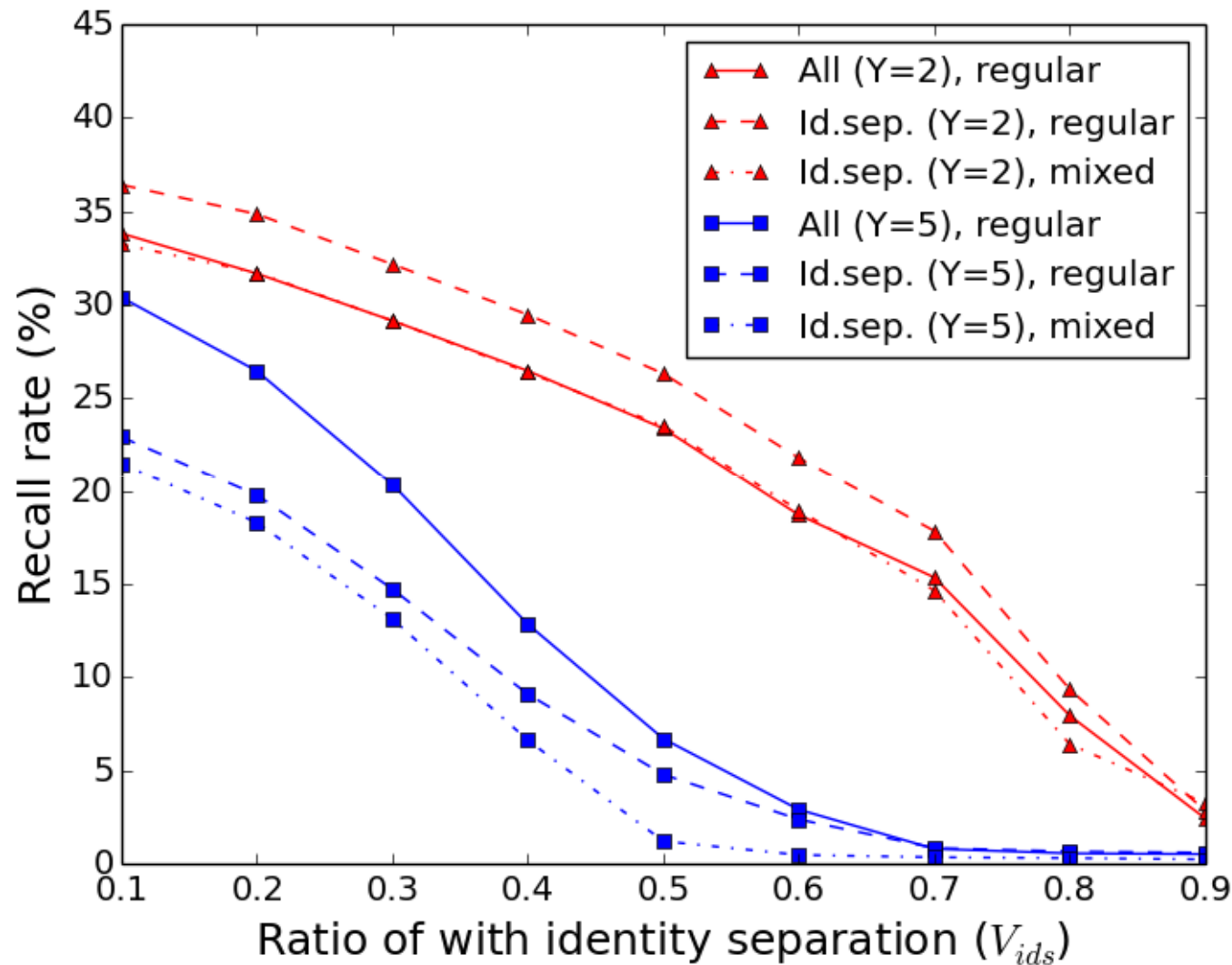
Feasible strategies!

[T2.2-3] Can identity separation provide protection for the network?

- I evaluated several settings of identity separation with structural de-anonymization attacks, e.g.:
 - Only using many separated identities is not enough (fig. on the left)
 - Plus deleting some edges introduces individual privacy only (fig. on the right)

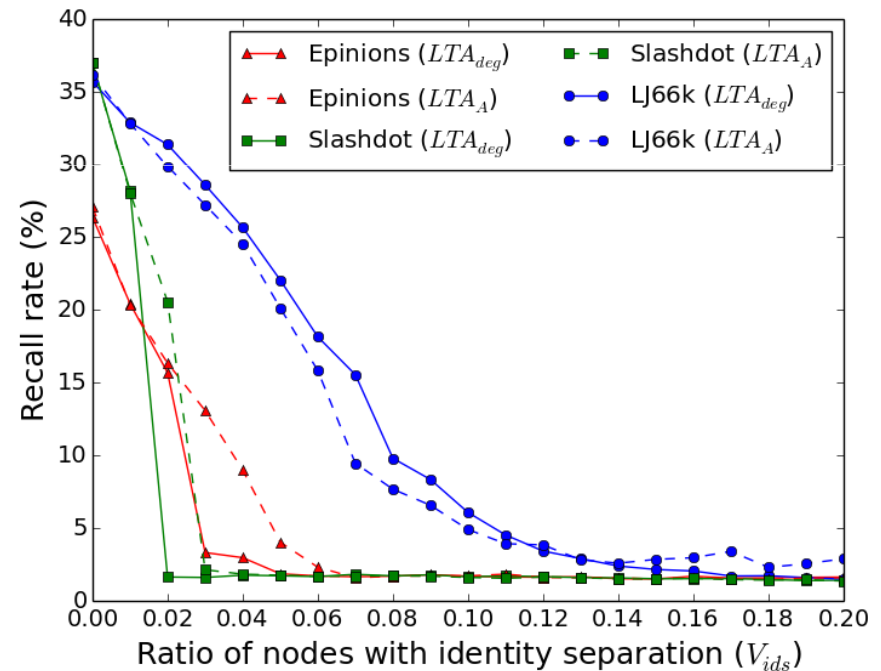
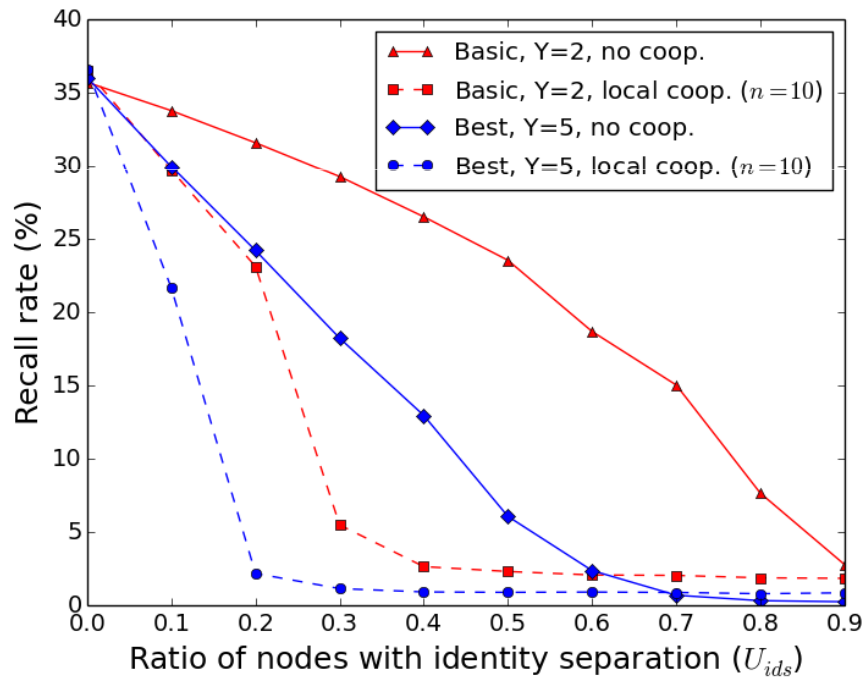


[T2.2-3] Can identity separation provide protection for the network? (2)



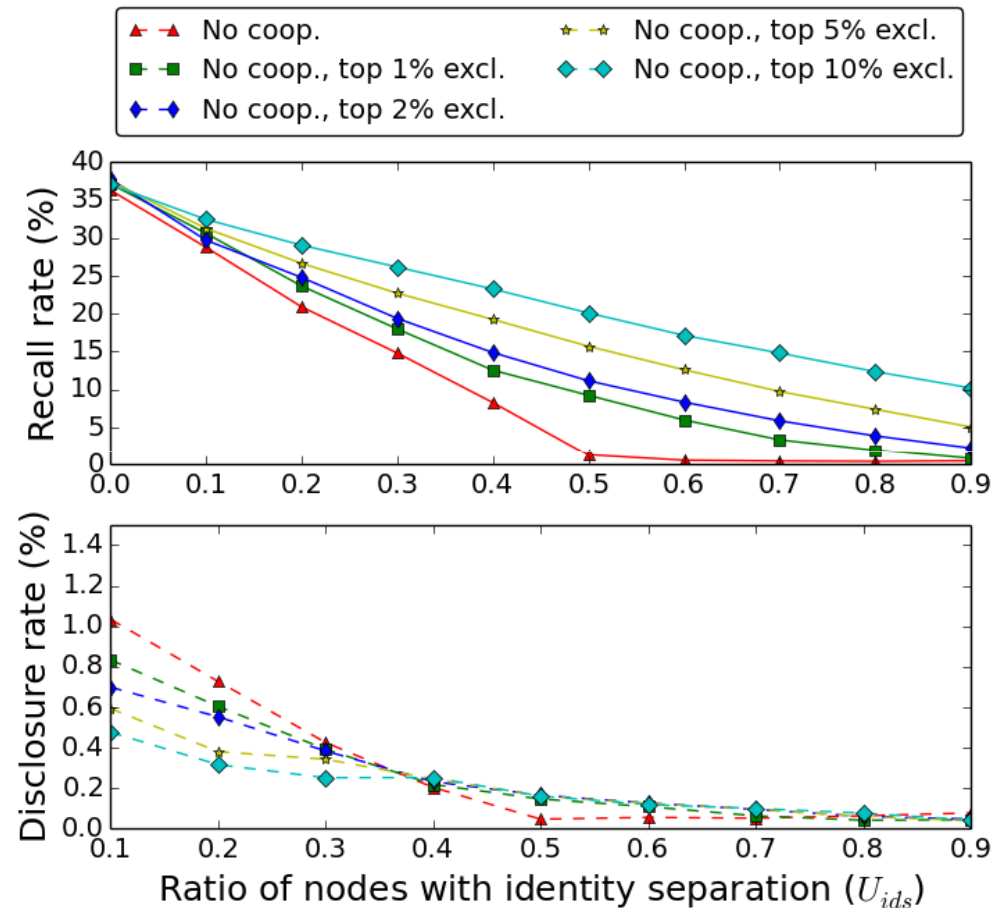
[T2.4-5] Can cooperation protect the network?

- Evaluation of different cooperation models:
 - Locally organized cooperation
 - Global cooperation: important nodes (low anonymity) should act



[T2.6] What happens if top degree nodes do not cooperate?

- Without their support, the performance of protection of network privacy degrades rapidly.

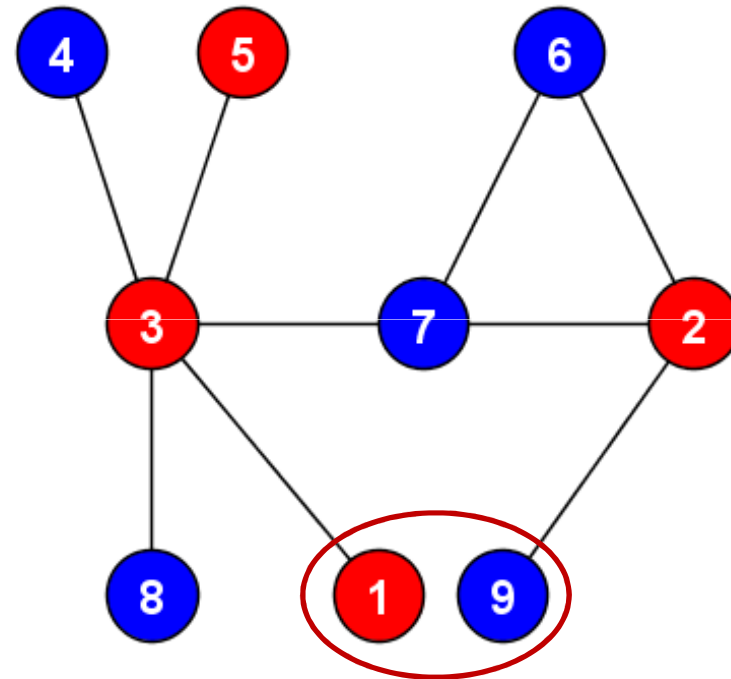


Related publications to Problem Set 2

- [J1] G. G. Gulyás and S. Imre, “Hiding information against structural re-identification,” *Telecommunication Systems*, September 2014. (under review).
- [J2] B. Simon, G. G. Gulyás, and S. Imre, “Analysis of grasshopper, a novel social network de-anonymization algorithm,” *Periodica Polytechnica Electrical Engineering and Computer Science*, January 2015. (accepted for publication).
- [J3] G. G. Gulyás and S. Imre, “Using identity separation against de-anonymization of social networks,” *Transactions on Data Privacy*, January 2015. (accepted for publication).
- [J4] G. G. Gulyás and S. Imre, “Analysis of identity separation against a passive clique-based de-anonymization attack,” *Infocommunications Journal*, vol. 4, pp. 11–20, December 2011.
- [C2] G. G. Gulyás and S. Imre, “Hiding information in social networks from de-anonymization attacks by using identity separation,” in *Communications and Multimedia Security* (B. Decker, J. Dittmann, C. Kraetzer, and C. Vielhauer, eds.), vol. 8099 of *Lecture Notes in Computer Science*, Springer Berlin Heidelberg, 2013.
- [C7] G. G. Gulyás, R. Schulcz, and S. Imre, “Modeling role-based privacy in social networking services,” in *Third International Conference on Emerging Security Information, Systems and Technologies (SECURWARE 2009)*, pp. 173–178, June 2009.
- [C8] G. G. Gulyás, “Design of an anonymous instant messaging service,” in *Proceedings of PET Convention 2009.1* (S. Köpsell and K. Loesing, eds.), pp. 34–40, Fakultät Informatik, TU Dresden, March 2009.

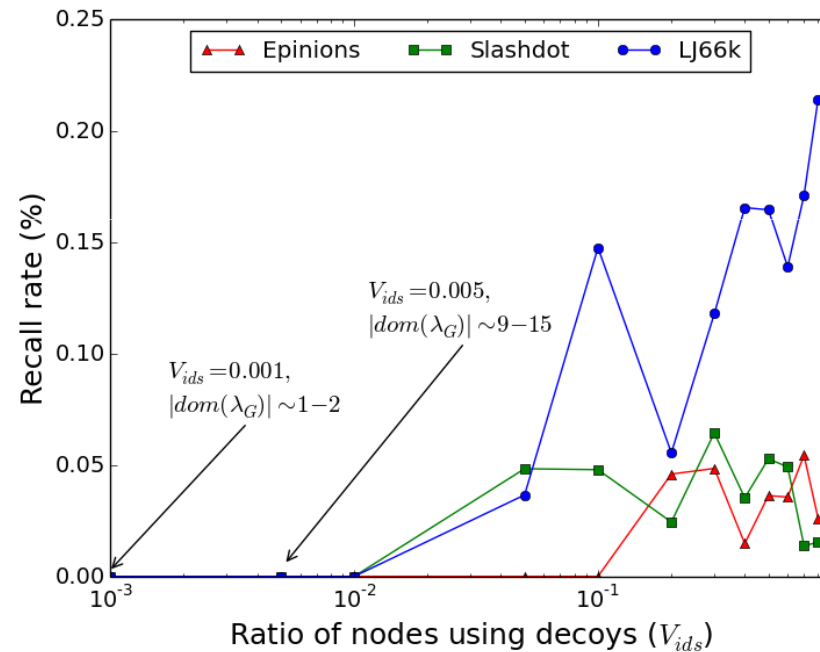
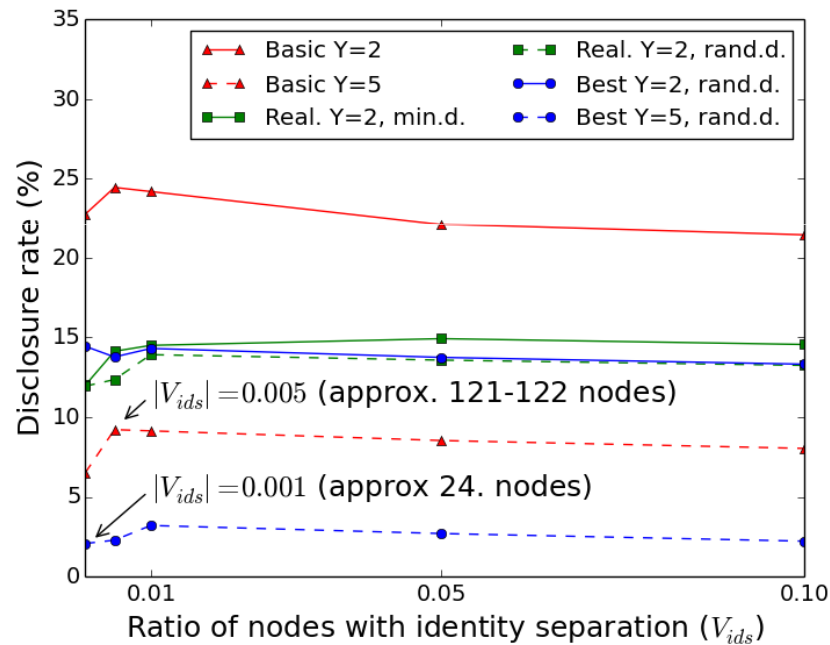
Problem Set 3: *individual strategies with identity separation*

- Are there any advised individual strategies?
 - decoying: the naive approach
 - (reversibility of identity separation)
 - adopting k-anonymity to the context
 - y-identity: a new model



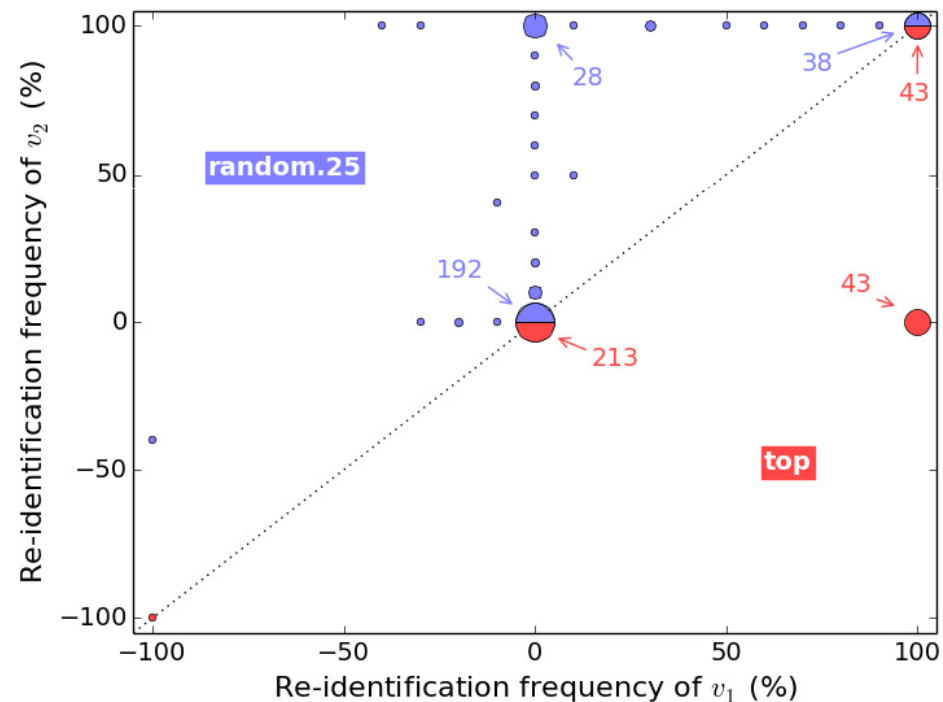
[T3.1] How effective a naive individual strategy can be?

- Handful of users are still good.
- Targeted information hiding: using decoy identities.



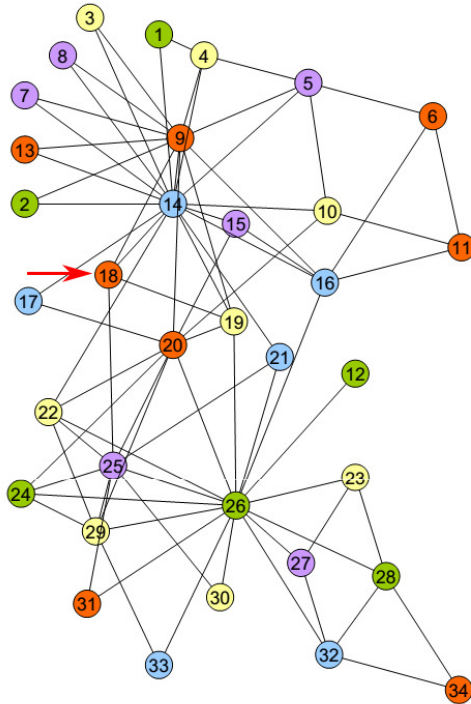
[T3.2] Can we reverse identity separation?

- I provided a method for calculating the lower bound of the probability of the discovery of partial identities with a simple modification of the state-of-the-art attack.

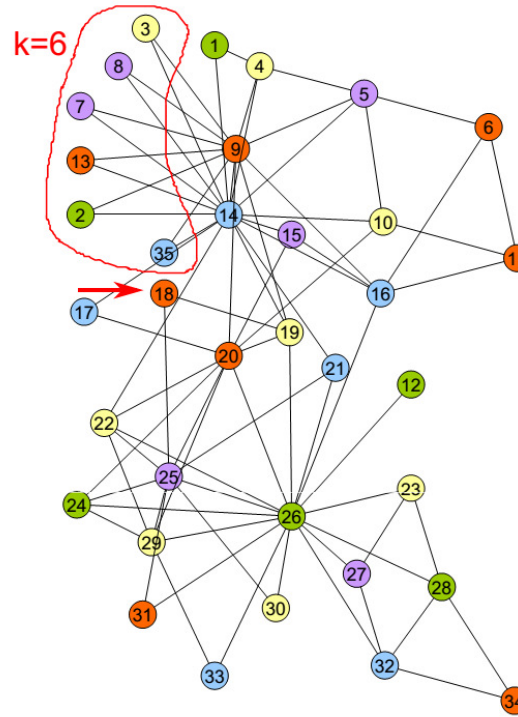


[T3.3] Can a k-anonymity variant enhance privacy?

Original network



k-anonymity example



Algorithm 1: $(k, 2)$ -anonymity with edge modification. It takes as input: the graph structure G , a node v_i selected for identity separation, c denoting the number of connections to anonymize, and parameter k of k-anonymity.

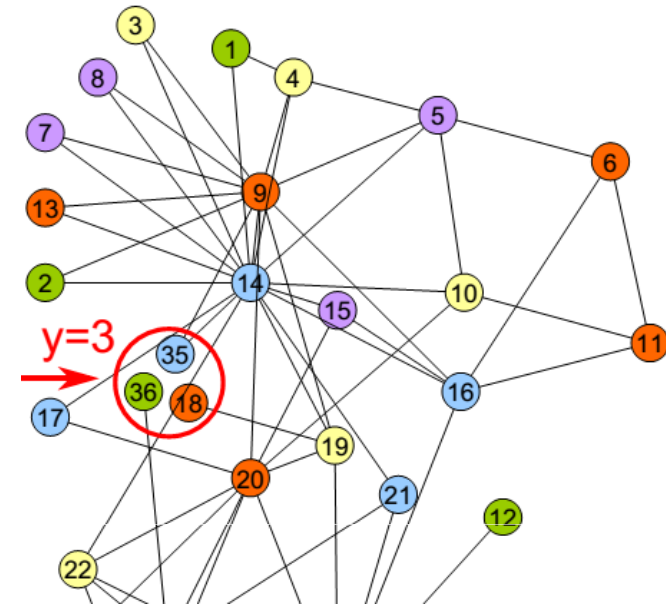
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1: procedure K-ANONYMIZENODE( $G, v_i, c, k$ )
2:   Calculate  $V_i, V_i^2$ 
3:    $c' \leftarrow c, V_k \leftarrow \{\}, E_k \leftarrow \{\}$ 
4:   while  $c' \geq 1$  and  $|V_k| = 0$  do
5:      $\kappa \leftarrow \{\}$  ▷ Groups having  $c'$  common neighbors with  $v_i$ 
6:     for all  $v_j \in V_i^2$  do
7:        $V_{\cap j} \leftarrow V_i \cap G.nbrs(v_j)$ 
8:       if  $|V_j| = c$  and  $|V_{\cap j}| = c'$  then
9:          $\kappa[V_{\cap j}] \leftarrow \kappa[V_{\cap j}] \cup \{v_j\}$ 
10:      end if
11:    end for
12:    for all  $\kappa[V_{\cap j}]$  if  $|\kappa[V_{\cap j}]| \geq k - 1$  do
13:      if  $c = c'$  then ▷ k-anonymity without modification
14:         $V_k \leftarrow \kappa[V_{\cap j}]$ 
15:        break
16:      end if
17:       $\psi \leftarrow \{\}$  ▷ Get new neighbors related to the k-group
18:      for all  $v_j \in \kappa[V_{\cap j}]$  do
19:         $V_{j \setminus i} \leftarrow G.nbrs(v_j) \setminus V_i \setminus \kappa[V_{\cap j}] \setminus \{v_i\}$ 
20:        for all  $v_t \in V_{j \setminus i}$  do
21:           $\psi[v_t] \leftarrow G.nbrs(v_t) \cap \kappa[V_{\cap j}]$ 
22:        end for
23:      end for
24:       $\eta \leftarrow \{\}$  ▷ Filter applicable groups and neighbors
25:      for all  $\psi[v_t]$  do
26:        for all  $\gamma \subseteq \psi[v_t]$  if  $|\gamma| = k - 1$  do
27:           $\eta[\gamma] \leftarrow \eta[\gamma] \cup \{v_t\}$ 
28:        end for
29:      end for
30:      if  $\exists \eta[\gamma]$  that  $|\eta[\gamma]| \geq c - c'$  then
31:        pick  $\eta[\gamma]$  where  $|\eta[\gamma]| \geq c - c'$ 
32:         $V_k \leftarrow \gamma$ 
33:         $E_k \leftarrow \eta[\gamma]$ 
34:        break
35:      end if
36:    end for
37:     $c' = c' - 1$ 
38:  end while
39:  return  $V_k, E_k$  ▷ Existing and new neighbors for k-anonymity
40: end procedure

```

[T3.4] Is there a theoretically best approach?

- The y -identity model
- Definition as a game:
 - Player set \mathcal{P} , Strategy set \mathcal{S} , utility set \mathcal{U}
 - Strong attackers know they revealed \mathcal{S}
 - Weak attackers reveal only $\mathcal{S}' \subseteq \mathcal{S}$ (with uncertainty)
- For the given attacker model
 - Proven best strategies for each attacker type
 - There is proven a feasible strategy against an unknown attackers



y -identity model

Related publications to Problem Set 3

[J1] G. G. Gulyás and S. Imre, “Hiding information against structural re-identification,” *Telecommunication Systems*, September 2014. (under review).

[J3] G. G. Gulyás and S. Imre, “Using identity separation against de-anonymization of social networks,” *Transactions on Data Privacy*, January 2015. (accepted for publication).

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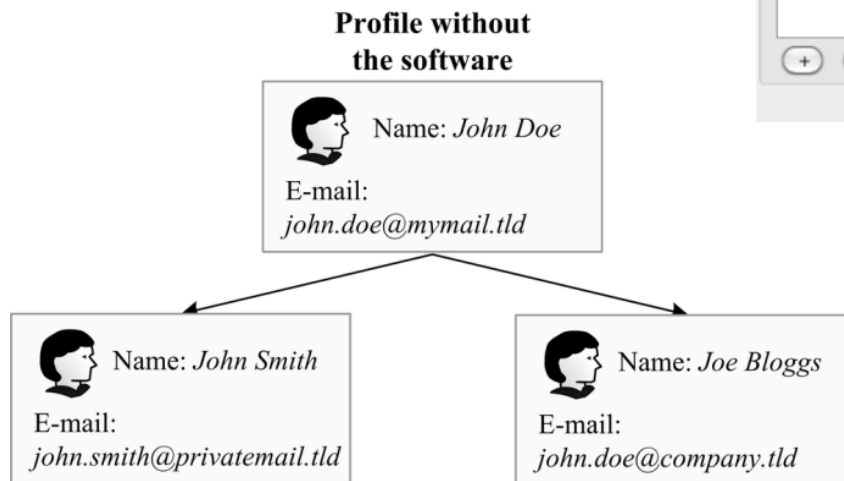
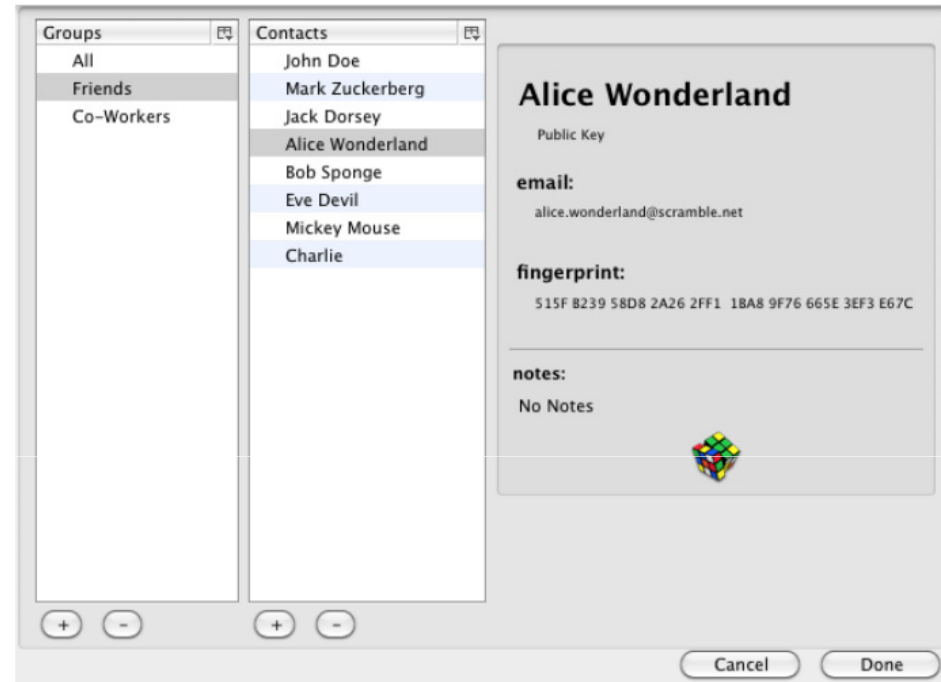
[C7] G. G. Gulyás, R. Schulcz, and S. Imre, “Modeling role-based privacy in social networking services,” in *Third International Conference on Emerging Security Information, Systems and Technologies (SECURWARE 2009)*, pp. 173–178, June 2009.

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Applications of Results

F. Beato, et al. "Scramble! your social network data.", 2011.

- Conclusion
 - Global protection is hard
 - Individually it works
- Client side implementation?
 - Strategies similarly as proposed



Publications statistics

Type	Hungaran	National	Internat.	Count	Score
Journal					
Peer-reviewed		2	1	3	12
Non peer-reviewed	1			1	½
Book					
Book chapter	4	1	2	7	21
Conference papers			9	9	18
Patents					
Others	2			2	5/6
Sum	7	4	11	22	52.3

- 3 peer-reviewed international journal papers (all in Scopus)
 - +1 currently under review in a WoS journal (only minor revision)
- 42+4 independent citations

Acknowledgement goes to...

- Sándor Imre
- Levente Buttyán
- Iván Székely
- Colleagues for inspiring discussions: Tamás Holczer, Márk Félegyházi, Máté Horváth, András Telcs
- Co-authors
- Members of the CrySyS & MCL Labs
- Reviewers of my dissertation: Julien Freudiger and Gergely Biczók
- My Wife, and ...

Thank you for your attention!
Any questions?



Gábor György Gulyás

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REVIEWER QUESTIONS

Julien Freudiger #1

- *What is the best way to bring the multiple identity management to market? Should we assist users into managing multiple identities, or should social networks adopt such practices?*
- Possible targets: 'natural' users of identity separation and privacy conscious folks
- Client/server side: depends on the service
 - Client side: web identity management, MAC change (+social netw.)
 - Server side: social networks
- Some apps/features already show that the idea could work, e.g., hmmmapp.com or circles in Google+

Gergely Biczók #1

- *The candidate has tapped into an intriguing line of research with his work on simple identity partitioning games in Chapter 6. Which family of (naturally more complex) game-theoretical models could be utilized to better capture the connected nature of the social networks under attack?*
- Possible extensions are mentioned in my dissertation, e.g., another type of weak attacker who can assess the possible unknown strategy set of the user
- Connections could be according to the model in [1] to evaluate how the adoption of the identity separation. This model focuses on the cost-benefit ratio.

[1] Ohtsuki, Hisashi, et al. "A simple rule for the evolution of cooperation on graphs and social networks." *Nature* 441.7092 (2006): 502-505.